

# ARITHMETICAL FUNCTIONALS AND WEAK KÖNIG'S LEMMA

KAZUYUKI TANAKA AND TAKESHI YAMAZAKI

ABSTRACT. By  $\text{RCA}_0$ , we denote the system of recursive comprehension axioms with  $\Sigma_1^0$  induction.  $\text{WKL}_0$  is  $\text{RCA}_0$  plus weak König's lemma: every infinite tree of sequences of 0's and 1's has an infinite path. In this paper, we first show that for any countable model  $(M, S)$  of  $\text{RCA}_0$ , there exists a countable model  $(M, S')$  of  $\text{WKL}_0$  such that  $S \cap S'$  consists of all  $\Delta_1^0$  subsets of  $M$ . By sophisticating this argument, we finally prove that if a functional is arithmetically definable in  $\text{WKL}_0$ , it is already so in  $\text{RCA}_0$ . More precisely, it is shown that if  $\text{WKL}_0$  proves  $\forall x \forall X \exists! Y \varphi(x, X, Y)$  with  $\varphi$  arithmetical, so does  $\text{RCA}_0$ .

## 1. INTRODUCTION

A celebrating metamathematical theorem due to L. Harrington asserts that  $\text{WKL}_0$  is conservative over  $\text{RCA}_0$  for the arithmetical (in fact,  $\Pi_1^1$ ) sentences. In other words, if an arithmetical theorem can be obtained by some analytical methods involving the compactness argument over computable mathematics, it is already provable without it. This result can be viewed as a computable analogue of the Gödel-Kreisel theorem on set theory, which asserts that if an arithmetical sentence can be proved in ZF with the axiom of choice, it is already provable without it.

Then it is natural to think of extending Harrington's conservation result to analytical sentences, since the Gödel-Kreisel theorem has been extended to the  $\Sigma_2^1$  sentences by J. Shoenfield. However, it is obvious that  $\text{WKL}_0$  is not  $\Sigma_1^1$  conservative over  $\text{RCA}_0$ , since an instance of weak König's lemma is  $\Sigma_1^1$ .

In this context, we claim that if  $\text{WKL}_0$  proves  $\forall x \forall X \exists! Y \varphi(x, X, Y)$  with  $\varphi$  arithmetical, so does  $\text{RCA}_0$ . Note that  $\exists! X \varphi(X)$  means that there exists a unique  $X$  satisfying  $\varphi(X)$ . This claim answers a problem posed by the first author [9] and attempted by some others [2], [10].

Let us see an application of our claim at first. The fundamental theorem of algebra, which asserts that any complex polynomial of any positive degree has a zero (or a unique non-void set of roots), can be stated in the form  $\forall x \forall X \exists! Y \varphi(x, X, Y)$  by using a canonical expression (i.e., the binary expansion) for the complex numbers. Most of popular proofs of the theorem use some analytical methods which can be easily formalized in  $\text{WKL}_0$  but not in  $\text{RCA}_0$ . However, by our conservation result, it can be concluded without elaborating a computable solution that the fundamental theorem of algebra is already provable in  $\text{RCA}_0$ .

By contrast, take a look at the statement that any continuous real function on the closed unit interval  $[0, 1]$  has a maximum value. This sentence can not be expressed in the form  $\forall x \forall X \exists! Y \varphi(x, X, Y)$  with  $\varphi$  arithmetical. The point is that

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we can not determine arithmetically whether or not a set encodes a *total* continuous function in the terms of Simpson [8].

Now, we give rigorous definitions of the systems  $\text{RCA}_0$  and  $\text{WKL}_0$ . The language  $\mathcal{L}_2$  of second-order arithmetic is a two-sorted language with number variables  $x, y, z, \dots$  and set variables  $X, Y, Z, \dots$ . Numerical terms are built up from numerical variables and constant symbols  $0, 1$  by means of binary operations  $+$  and  $\cdot$ . Atomic formulas are  $s = t$ ,  $s < t$  and  $s \in X$ , where  $s$  and  $t$  are numerical terms. *Bounded* ( $\Sigma_0^0$  or  $\Pi_0^0$ ) formulas are constructed from atomic formulas by propositional connectives and bounded numerical quantifiers ( $\forall x < t$ ) and ( $\exists x < t$ ), where  $t$  does not contain  $x$ . A  $\Sigma_n^0$  formula is of the form  $\exists x_1 \forall x_2 \dots x_n \theta$  with  $\theta$  bounded, and a  $\Pi_n^0$  formula is of the form  $\forall x_1 \exists x_2 \dots x_n \theta$  with  $\theta$  bounded. All the  $\Sigma_n^0$  and  $\Pi_n^0$  formulas are the *arithmetical* ( $\Sigma_1^1$  or  $\Pi_1^1$ ) formulas. A  $\Sigma_n^1$  formula is of the form  $\exists X_1 \forall X_2 \dots X_n \varphi$  with  $\varphi$  arithmetical, and a  $\Pi_n^1$  formula is of the form  $\forall X_1 \exists X_2 \dots X_n \varphi$  with  $\varphi$  arithmetical.

The system of  $\text{RCA}_0$  consists of

1. the ordered semiring axioms for  $(\omega, +, \cdot, 0, 1, <)$ ,
2.  $\Delta_1^0$  comprehension scheme:

$$\forall x(\varphi(x) \leftrightarrow \psi(x)) \rightarrow \exists X \forall x(x \in X \leftrightarrow \varphi(x)),$$

where  $\varphi(x)$  is  $\Sigma_1^0$ ,  $\psi(x)$  is  $\Pi_1^0$ , and  $X$  does not occur freely in  $\varphi(x)$ ,

3.  $\Sigma_1^0$  induction scheme:

$$\varphi(0) \wedge \forall x(\varphi(x) \rightarrow \phi(x+1)) \rightarrow \forall x \varphi(x),$$

where  $\varphi(x)$  is a  $\Sigma_1^0$  formula.

Within  $\text{RCA}_0$ , we define  $2^{<\mathbb{N}}$  to be the set of (codes for) finite sequences of 0's and 1's. A set  $T \subseteq 2^{<\mathbb{N}}$  is said to be a *tree* (or precisely 0-1 *tree*) if any initial segment of a sequence in  $T$  is also in  $T$ . A *path* through  $T$  is a function  $f : \mathbb{N} \rightarrow \{0, 1\}$  such that for each  $n$ , the sequence  $f[n] = \langle f(0), f(1), \dots, f(n-1) \rangle$  belongs to  $T$ . The axioms of  $\text{WKL}_0$  consists of those of  $\text{RCA}_0$  plus *weak König's lemma*: every infinite 0-1 tree  $T$  has a path.

The interest of  $\text{WKL}_0$  has been well established through an ongoing program, called *Reverse Mathematics*. H. Friedman, S. G. Simpson and others have shown that numerous well-known theorems in different fields of mathematics are provably equivalent to  $\text{WKL}_0$  over  $\text{RCA}_0$  [8].

An  $\mathcal{L}_2$ -structure  $(M, S, +_M, \cdot_M, 0_M, 1_M, <_M, \in_{M \times S})$  is simply denoted by  $(M, S)$  throughout this paper. We also write  $M$  for an  $\mathcal{L}_1$  structure  $(M, +_M, \cdot_M, 0_M, 1_M, <_M)$ . If  $M$  is the set (or structure) of standard natural numbers, an  $\mathcal{L}_2$ -structure  $(M, S)$  is called an  $\omega$ -*structure* or an  $\omega$ -*model*. In particular,  $\omega$ -models of  $\text{WKL}_0$  are known as Scott systems and extensively studied by not a few people, e.g. Kaye [4].

In the next section, we prove that for any countable model  $(M, S)$  of  $\text{RCA}_0$ , there exists a countable model  $(M, S')$  of  $\text{WKL}_0$  such that  $S \cap S'$  is the set of  $\Delta_1^0$  subsets of  $M$ . This can be regarded as a non- $\omega$  extension of Kreisel's recursive hard core theorem, which asserts that the intersection of all  $\omega$ -models of  $\text{WKL}_0$  is the set of recursive sets.

In section 3, we refine the argument for the non- $\omega$  hard core theorem to show that for any countable model  $M$  of the axioms of ordered semirings and  $\Sigma_1^0$  induction, there exist two countable models  $(M, S)$  and  $(M, S')$  of  $\text{WKL}_0$  such that  $S \cap S'$  is the set of  $\Delta_1^0$  subsets of  $M$  and in addition,  $(M, S \cup S') \models \Sigma_1^0$  induction. The last condition implies that there exists a third model  $(M, S'')$  of  $\text{WKL}_0$  such that

$S \cup S' \subseteq S''$ . To see how to use this fact, let us suppose that  $\text{WKL}_0$  proves a sentence  $\exists! X \varphi(X)$  with  $\varphi$  arithmetical. Then both  $(M, S)$  and  $(M, S')$  have a unique set  $X$  satisfying  $\varphi(X)$ , and such unique sets must be identical (hence in  $\Delta_1^0$ ), for otherwise  $(M, S'')$  has more than one set satisfying  $\varphi$ . Therefore,  $\text{RCA}_0$  also proves  $\exists! X \varphi(X)$ .

## 2. A NON- $\omega$ HARD CORE THEOREM

In this section, we first review the tree forcing argument originated by Jockusch-Soare [3] and used by L.Harrington for his conservation result on  $\text{WKL}_0$ . We then reinforce this argument with some other machinery to prove that for any countable model  $(M, S)$  of  $\text{RCA}_0$ , there exists a countable model  $(M, S')$  of  $\text{WKL}_0$  such that  $S \cap S'$  is the set of  $\Delta_1^0$  subsets of  $M$ .

Let  $(M, S)$  be an  $\mathcal{L}_2$ -structure. We say that  $X \subseteq M$  is  $\Delta_1^0$  definable over  $(M, S)$ , denoted  $X \in \Delta_1^0(S)$ , if there exist a  $\Sigma_1^0$  formula  $\varphi$  and a  $\Pi_1^0$  formula  $\psi$  both with parameters from  $M \cup S$  such that

$$X = \{m \in M : (M, S) \models \varphi(m)\} = \{m \in M : (M, S) \models \psi(m)\}.$$

We write  $\Delta_1^0$  for  $\Delta_1^0(\emptyset)$ . It is easy to see that if  $(M, S)$  is a model of  $\text{RCA}_0$ ,  $\Delta_1^0(S) = S$ .

**Lemma 1.** *Let  $(M, S)$  be an  $\mathcal{L}_2$ -structure which satisfies the axioms of ordered semirings and  $\Sigma_1^0$  induction. Then  $(M, \Delta_1^0(S))$  is a model of  $\text{RCA}_0$ .*

*Proof.* See the proof of Lemma IX.1.8 [8].  $\square$

We now define basic notions of the tree forcing. Let  $(M, S)$  be a model of  $\text{RCA}_0$ . Let  $\mathcal{T}_S$  be the set of all  $T \in S$  such that

$$(M, S) \models T \text{ is an infinite 0-1 tree with no path.}$$

Note that  $\mathcal{T}_S = \emptyset$  if and only if  $(M, S)$  is a model of  $\text{WKL}_0$ . We say that  $D \subseteq \mathcal{T}_S$  is *dense* if for each  $T \in \mathcal{T}_S$ , there exists  $T' \in D$  such that  $T' \subseteq T$ . For a set  $\mathcal{D}$  of dense sets,  $P$  is said to be *generic* for  $\mathcal{D}$  if for each set  $D \in \mathcal{D}$ , there exists  $T \in D$  such that  $P$  is a path through  $T$ .

**Lemma 2.** *Let  $(M, S)$  be a countable model of  $\text{RCA}_0$ , and  $\mathcal{D}$  a countable set of dense subsets of  $\mathcal{T}_S$ . Then each  $T \in \mathcal{T}_S$  has a path generic for  $\mathcal{D}$ .*

*Proof.* If  $\mathcal{D} = \{D_i : i \in \omega\}$  is a set of dense sets, we can easily construct a sequence of trees  $T_i$  ( $i \in \omega$ ) such that  $T_0 = T, T_{i+1} \subseteq T_i$  and  $T_{i+1} \in D_i$  for each  $i \in \omega$ . Then a path  $P \subseteq \bigcap T_i$  is what we want.  $\square$

**Lemma 3.** *Let  $(M, S)$  be a model of  $\text{RCA}_0$ . For any tree  $T \in \mathcal{T}_S$ , there exists a path  $P$  through  $T$  such that  $(M, S \cup \{P\}) \models \Sigma_1^0$  induction.*

*Proof.* Let  $(M, S)$  be a model of  $\text{RCA}_0$ . Since  $\Sigma_1^0$  induction is provably equivalent to bounded  $\Sigma_1^0$  comprehension (cf. Remark II.3.11 [8]), it suffices to prove that any tree  $T \in \mathcal{T}_S$  has a path  $P$  such that for each  $m \in M$ ,  $\{n \in M : n <_M m \wedge (M, S \cup \{P\}) \models \varphi(n, P)\} \in S$ , where  $\varphi(x, X)$  is a  $\Sigma_1^0$  formula with parameters from  $M \cup S$ .

Let  $\{\varphi_e(x, X) : e \in \omega\}$  enumerate all the  $\Sigma_1^0$ -formulas with parameters from  $M \cup S$  and no free variables other than  $x$  and  $X$ . Without loss of generality, we may assume that  $\varphi_e(x, X)$  is of the form  $\exists y \theta_e(x, X[y])$  with  $\theta_e(x, s) \in \Sigma_0^0$ , where  $X[y]$  denotes the sequence  $\langle f(0), f(1), \dots, f(y-1) \rangle$  with the characteristic function  $f$  of  $X$ .

For each  $e \in \omega$  and  $m \in M$ , let  $D_{e,m}^0$  be the set of all  $T \in \mathcal{T}_S$  such that for any  $n <_M m$ ,  $(M, S)$  satisfies either

1.  $\forall s \in T \neg \theta_e(n, s)$ , or
2.  $\exists w \forall s \in T (\text{lh}(s) = w \rightarrow \exists y \leq w \theta_e(n, s[y]))$ ,

where  $\text{lh}(s)$  denotes the length of sequence  $s$ , and  $s[y]$  is the initial subsequence of  $s$  with the length  $y$ . Then it is not difficult to see that  $D_{e,m}^0$ 's are dense. (See Lemma IX.2.4 [8].)

Let  $T \in \mathcal{T}_S$  be given. By Lemma 2, we can take a path  $P$  through  $T$  which is generic for  $\{D_{e,m}^0 : e \in \omega, m \in M\}$ . Fix any  $e \in \omega$  and  $m \in M$ . Since  $P$  is generic, there is a tree  $T' \in D_{e,m}^0$  which has a path  $P$ . Then, it is easy to see

$$\begin{aligned} & \{n \in M : n <_M m \wedge (M, S \cup \{P\}) \models \varphi_e(n, P)\} \\ &= \{n \in M : n <_M m \wedge (M, S) \models \exists w \forall s \in T' (\text{lh}(s) = w \rightarrow \exists y \leq w \theta_e(n, s[y]))\}. \end{aligned}$$

The set on the right-hand side belongs to  $S$  by bounded  $\Sigma_1^0$  comprehension for  $(M, S)$ . Thus  $(M, S \cup \{P\})$  also satisfies bounded  $\Sigma_1^0$  comprehension.  $\square$

**Lemma 4.** *Let  $(M, S)$  be a countable model of  $\text{RCA}_0$ . For any infinite 0-1 tree  $T \in S$ , there exists a countable model  $(M, S')$  of  $\text{RCA}_0$  such that  $S \subseteq S'$  and  $T$  has a path in  $S'$ .*

*Proof.* It is obvious from Lemmas 1 and 3.  $\square$

**Lemma 5.** *Let  $(M, S)$  be a countable model of  $\text{RCA}_0$ . Then there exists a countable model  $(M, S')$  of  $\text{WKL}_0$  such that  $S \subseteq S'$ .*

*Proof.* Use Lemma 4 repeatedly.  $\square$

**Theorem 6.** (*L. Harrington*) *For any  $\Pi_1^1$ -sentence  $\varphi$ , if  $\varphi$  is a theorem of  $\text{WKL}_0$ , then  $\varphi$  is already a theorem of  $\text{RCA}_0$ . Especially, the first order part of  $\text{WKL}_0$  is the same as that of  $\text{RCA}_0$ , i.e.,  $\text{I}\Sigma_1$  (Peano arithmetic with induction restricted to the  $\Sigma_1^0$   $\mathcal{L}_1$ -formulas).*

*Proof.* It easily follows Lemma 5 by the help of Gödel's completeness theorem.  $\square$

We now recall another important characterization of models of  $\text{WKL}_0$ .

**Theorem 7.** *There is a  $\Pi_1^0$ -formula  $\psi(X, Y)$  with no free variables other than  $X$  and  $Y$  such that for any model  $(M, S)$  of  $\text{WKL}_0$  and for any  $A \in S$ ,*

- (1) *there exists  $W \in S$  such that  $(M, S) \models \psi(A, W)$ , and*
- (2) *if  $(M, S) \models \psi(A, W)$ , then  $(M, \{(W)_n : n \in M\}) \models \text{WKL}_0$  and  $A \in \{(W)_n : n \in M\}$ , where  $(W)_n = \{k \in M : (k, n) \in W\}$  for each  $n \in M$ .*

*Proof.* See Lemma VII.2.9 [8].  $\square$

The above theorem essentially says that  $\text{WKL}_0$  proves the existence of a structure satisfying  $\text{WKL}_0$ . We notice that this assertion does not conflict with Gödel's second incompleteness theorem, since the structure need not be equipped with the satisfaction relation. Though the essence of the theorem is a kind of folklore, the above particular statement is due to S. Simpson [8]. See also [4], [5] for other accounts.

Since  $\{(W)_n : n \in M\}$  is a proper subset of  $S$  in the theorem, it turns out that there is no minimal model of  $\text{WKL}_0$  with a fixed first order part  $M$ . Later in this section, we will show that for a fixed  $M$ , the intersection of all such  $S$ 's that  $(M, S) \models \text{WKL}_0$  is just  $\Delta_1^0$ . If  $M = \omega$ , this fact has been known as Kreisel's hard core theorem.

**Lemma 8.** *Let  $(M, S)$  be a countable model of  $\text{RCA}_0$ . Let  $T$  be a tree in  $\mathcal{T}_S$ . Then, for each  $A \subseteq M$  such that  $A \notin S$ , there exists a path  $P$  through  $T$  such that  $A$  is not in  $\Delta_1^0(S \cup \{P\})$  and that  $(M, S \cup \{P\}) \models \Sigma_1^0$  induction.*

*Proof.* Let  $(M, S)$  be a countable model of  $\text{RCA}_0$  and  $A$  a subset of  $M$  such that  $A \notin S$ . Let  $\{\varphi_e(x, X) : e \in \omega\}$  enumerates all the  $\Sigma_1^0$ -formulas with parameters from  $M \cup S$  and no free variables other than  $x$  and  $X$ .

We first claim that for each  $T \in \mathcal{T}_S$  and each pair  $(e, d) \in \omega^2$ , there exists a path  $Z$  through  $T$  such that  $(M, \{Z\}) \models \Sigma_1^0$  induction and that  $(e, d)$  is not a  $\Delta_1^0(S \cup \{Z\})$ -index of  $A$ , i.e.,  $A \neq \{m : (M, S \cup \{Z\}) \models \varphi_e(m, Z)\}$  or  $A \neq \{m : (M, S \cup \{Z\}) \models \neg\varphi_d(m, Z)\}$ .

By way of contradiction, deny the claim. Then there exist a tree  $T \in \mathcal{T}_S$  and a pair  $(e, d) \in \omega^2$  such that if  $Z$  is a path through  $T$  and  $(M, S \cup \{Z\}) \models \Sigma_1^0$  induction, then  $(e, d)$  is a  $\Delta_1^0(S \cup \{Z\})$ -index of  $A$ . By Lemma 5, we can construct a countable model  $(M, S')$  of  $\text{WKL}_0$  such that  $S \subseteq S'$ . Then, for any path  $Z$  through  $T$  such that  $Z \in S'$ ,

$$\begin{aligned} m \in A &\Leftrightarrow (M, S \cup \{Z\}) \models \varphi_e(m, Z) \\ &\Leftrightarrow (M, S') \models \varphi_e(m, Z). \end{aligned}$$

Hence, we have

$$m \in A \Leftrightarrow (M, S') \models \forall Z [Z \text{ is a path through } T \rightarrow \varphi_e(m, Z)].$$

Since “ $Z$  is a path through  $T$ ” is expressed as a  $\Pi_1^0$  formula, “ $Z$  is a path through  $T \rightarrow \varphi_e(m, Z)$ ” is  $\Sigma_1^0$ , and so the whole formula  $\forall Z [Z \text{ is a path through } T \rightarrow \varphi_e(m, Z)]$  is logically equivalent in  $(M, S')$  to a  $\Sigma_1^0$  formula  $\varphi'(m)$  (with parameters from  $M \cup S$ ) by virtue of compactness of the Cantor space (cf. Lemma V.III.2.4 [8]). Since for any  $m \in M$ ,  $(M, S') \models \varphi'(m)$  if and only if  $(M, S) \models \varphi'(m)$ , we finally have

$$m \in A \Leftrightarrow (M, S) \models \varphi'(m).$$

Similarly, we have

$$m \in A \Leftrightarrow (M, S') \models \exists Z [Z \text{ is a path through } T \wedge \neg\varphi_d(m, Z)],$$

and so by compactness, there exists a  $\Pi_1^0$ -formula  $\psi'(m)$  with parameters from  $M \cup S$  such that

$$m \in A \Leftrightarrow (M, S) \models \psi'(m).$$

Therefore,  $A$  is in  $\Delta_1^0(S)$ , hence in  $S$  since  $(M, S)$  is a model of  $\text{RCA}_0$ . This contradicts with our assumption. Thus the claim is proved.

From now, we may assume that for each  $e \in \omega$ ,  $\Sigma_1^0$ -formula  $\varphi_e(x, X)$  is of the form  $\exists y \theta_e(x, X[y])$  with  $\theta_e(x, s) \in \Sigma_0^0$ . For each  $(e, d) \in \omega^2$ , we define  $D_{e,d}^A$  to be the set of all  $T \in \mathcal{T}_S$  such that one of the followings holds for some  $m \in M$ :

- A1.  $m \in A \wedge (M, S) \models \forall s \in T \neg\theta_e(m, s)$ ,
- A2.  $m \notin A \wedge (M, S) \models \exists w \forall s \in T (\text{lh}(s) = w \rightarrow \exists y \leq w \theta_e(m, s[y]))$ ,
- A3.  $m \in A \wedge (M, S) \models \exists w \forall s \in T (\text{lh}(s) = w \rightarrow \exists y \leq w \theta_d(m, s[y]))$ ,
- A4.  $m \notin A \wedge (M, S) \models \forall s \in T \neg\theta_d(m, s)$ .

To show that for each  $(e, d) \in \omega^2$ ,  $D_{e,d}^A$  is dense, we choose any  $T \in \mathcal{T}_S$ . By the above claim, there exists a path  $Z$  through  $T$  and  $m \in M$  such that one of the following conditions holds:

- B1.  $m \in A \wedge (M, S \cup \{Z\}) \models \forall y \neg\theta_e(m, Z[y])$ ,
- B2.  $m \notin A \wedge (M, S \cup \{Z\}) \models \exists y \theta_e(m, Z[y])$ ,
- B3.  $m \in A \wedge (M, S \cup \{Z\}) \models \exists y \theta_d(m, Z[y])$ ,

B4.  $m \notin A \wedge (M, S \cup \{Z\}) \models \forall y \neg \theta_d(m, Z[y])$ .

First suppose that condition B1 holds. Let  $T' = \{s \in T : \forall t \subseteq s \neg \theta_e(x, t)\}$ . Then,  $T' \in \mathcal{T}_S$ , since  $T' \in S$  and  $T'$  is an infinite subtree of  $T$ . It is also clear that A1 holds with  $T'$  (instead of  $T$ ). Thus  $T' \in D_{e,d}^A$ . Next suppose that condition B2 holds. Take an initial segment  $t$  of  $Z$  such that  $\theta_e(m, t)$ . Then  $T' = \{s \in T : s \subseteq t \text{ or } t \subseteq s\}$  clearly satisfies A2, hence  $T' \in D_{e,d}^A$ . The other two cases can be treated similarly. Hence, in any case, there exists a subtree  $T'$  of  $T$  such that  $T' \in D_{e,d}^A$ , which means that  $D_{e,d}^A$  is dense.

Given a  $T \in \mathcal{T}_S$ , we take a path  $P$  through  $T$  which is generic for  $\{D_{e,m}^0 : e \in \omega, m \in M\} \cup \{D_{e,d}^A : (e, d) \in \omega^2\}$ , where  $D_{e,m}^0$ 's are the dense sets defined in the proof of Lemma 3. Then,  $(M, S \cup \{P\})$  satisfies  $\Sigma_1^0$  induction by the proof of Lemma 3. By way of contradiction, we assume that  $A$  is in  $\Delta_1^0(S \cup \{P\})$ , that is, there exist  $e$  and  $d$  such that

$$A = \{m : (M, S \cup \{P\}) \models \exists y \theta_e(x, P[y])\} = \{m : (M, S \cup \{P\}) \models \forall y \neg \theta_d(x, P[y])\}.$$

Since  $P$  is generic, there exists  $T' \in D_{e,d}^A$  with a path  $P$ . First suppose that condition A1 of the definition of  $D_{e,d}^A$  holds for  $T'$ . Then there exists  $m \in A$  such that  $(M, S \cup \{P\}) \models \forall y \neg \theta_e(m, P[y])$ , since  $P \subset T'$ . This contradicts with the above equation for  $A$ . Suppose that condition A2 holds for  $T'$ . Then there exists  $m \notin A$  such that  $(M, S \cup \{P\}) \models \exists y \theta_e(m, P[y])$ , which is also absurd. Similarly, conditions A3 and A4 lead to a contradiction. Thus, we have shown that  $A$  is not in  $\Delta_1^0(S \cup \{P\})$ . This completes the proof.  $\square$

**Lemma 9.** *Let  $(M, S)$  be a countable model of  $\text{RCA}_0$ , and  $C$  a countable set of subsets of  $M$  such that  $C \cap S = \emptyset$ . Then any tree  $T \in \mathcal{T}_S$  has a path  $P$  such that  $C \cap \Delta_1^0(S \cup \{P\}) = \emptyset$  and that  $(M, S \cup \{P\}) \models \Sigma_1^0$  induction.*

*Proof.* Let  $(M, S)$  and  $C$  be as in the above statement. We define dense sets  $D_{e,m}^0$  and  $D_{e,d}^A$  as in the proofs of Lemmas 3 and 8. By Lemma 2, for each  $T \in \mathcal{T}_S$ , we can take a path  $P$  through  $T$  which is generic for  $\{D_{e,m}^0 : e \in \omega, m \in M\} \cup \{D_{e,d}^A : A \in C \wedge (e, d) \in \omega^2\}$ . Then by the proofs of Lemmas 3 and 8, it is easy to see that  $C \cap \Delta_1^0(S \cup \{P\}) = \emptyset$  and that  $(M, S \cup \{P\}) \models \Sigma_1^0$  induction.  $\square$

**Lemma 10.** *Let  $(M, S)$  be a countable model of  $\text{RCA}_0$ , and  $C$  a countable set of subsets of  $M$  such that  $C \cap S = \emptyset$ . Then there exists a countable model  $(M, S')$  of  $\text{WKL}_0$  such that  $S \subseteq S'$  and  $S' \cap C = \emptyset$ .*

*Proof.* Use the above lemma repeatedly.  $\square$

The next corollary is a generalized version of Kreisel's hard core theorem.

**Theorem 11.** *Let  $(M, S)$  be a countable model of  $\text{RCA}_0$ . Then there exists a countable model  $(M, S')$  of  $\text{WKL}_0$  such that  $S \cap S' = \Delta_1^0$ .*

*Proof.* By replacing  $S$  and  $C$  in Lemma 10 by  $\Delta_1^0$  and  $S - \Delta_1^0$ , respectively, we obtain the theorem.  $\square$

**Corollary 12.** *Let  $M$  be a countable model of  $\text{IS}_1$ . Then there exist uncountably many countable sets  $S$  of subsets of  $M$  such that  $(M, S)$  satisfies  $\text{WKL}_0$ .*

*Proof.* If there were only countably many of such  $S$ 's, then, putting  $C =$  (the union of all of them)  $-\Delta_1^0$ , by Lemma 10 we could obtain another model  $(M, S')$  of  $\text{WKL}_0$  such that  $S' \cap C = \emptyset$ , which is a contradiction.  $\square$

Finally, we remark that in our theorem,  $(M, S \cup S')$  may not satisfy  $\Sigma_1^0$  induction. In fact, there are two models  $(M, S)$  and  $(M, S')$  of  $\Sigma_1^0$  induction such that  $(M, S \cup S')$  does not satisfy  $\Sigma_1^0$  induction. This fact is easily obtained from the following two theorems.

**Theorem 13.** (*Mytilinaios* [7]) *Let  $M \models \mathbf{I}\Sigma_1$ , and let  $W$  be a  $\Sigma_1$  definable but not  $\Delta_1$  definable subset of  $M$ . Then there exist two  $\Sigma_1$  definable subsets  $A, B$  of  $M$  such that  $W$  is in  $\Delta_1^0(\mathbf{A}, \mathbf{B})$ , but neither in  $\Delta_1^0(\mathbf{A})$  nor in  $\Delta_1^0(\mathbf{B})$ .*

**Theorem 14.** (*Groszek et al.* [6]) *Let  $M \models \mathbf{B}\Sigma_2$ , and let  $A$  be a  $\Sigma_1$  definable subset of  $M$ . Then  $(M, \{A\}) \models \Sigma_1^0$  induction, or  $A$  is complete.*

Note that  $\mathbf{B}\Sigma_2$  denotes the collection axioms for the  $\Sigma_2$  formulas.

**Theorem 15.** *Let  $M \models \mathbf{B}\Sigma_2 \wedge \neg\mathbf{I}\Sigma_2$ . Then there exist two  $\Sigma_1$  definable subsets  $A, B$  of  $M$  such that both  $(M, \{A\})$  and  $(M, \{B\})$  satisfies  $\Sigma_1^0$  induction, but  $(M, \{A, B\})$  does not.*

*Proof.* Let  $M \models \mathbf{B}\Sigma_2 \wedge \neg\mathbf{I}\Sigma_2$ , and let  $W$  be a complete  $\Sigma_1$  subset of  $M$ . Then  $(M, \{W\})$  does not satisfy  $\Sigma_1^0$  induction, since  $M \models \neg\mathbf{I}\Sigma_2$  and any  $\Sigma_2$  set is  $\Sigma_1^0(W)$ . By Theorem 13, there exist two incomplete  $\Sigma_1$  subsets  $A, B$  of  $M$  such that  $W$  is in  $\Delta_1^0(\mathbf{A}, \mathbf{B})$ . By Theorem 14, both  $(M, \{A\})$  and  $(M, \{B\})$  satisfies  $\Sigma_1^0$  induction. But  $(M, \{A, B\})$  does not satisfy  $\Sigma_1^0$  induction, since  $(M, \{W\})$  does not.  $\square$

### 3. A NEW CONSERVATION RESULT

In this section, we prove our new conservation result by refining the preceding argument for the non- $\omega$  hard core theorem (Theorem 11). As a main lemma, we shall show that for any countable model  $M$  of the axioms of ordered semirings and  $\Sigma_1$  induction, there exist two countable models  $(M, S)$  and  $(M, S')$  of  $\mathbf{WKL}_0$  such that  $S \cap S'$  is the set of  $\Delta_1^0$  subsets of  $M$  and in addition,  $(M, S \cup S') \models \Sigma_1^0$  induction (Lemma 18).

Let  $(M, S)$  be a countable model of  $\mathbf{RCA}_0$ . Recall that  $\mathcal{T}_S$  denotes the set of all  $T \in S$  such that  $(M, S) \models T$  is an infinite 0-1 tree with no path. Let  $\mathcal{T}_S^2$  be the set of all pairs  $(T_0, T_1)$  of trees in  $\mathcal{T}_S$ . We define a partial order  $\leq$  on  $\mathcal{T}_S^2$  by  $(T'_0, T'_1) \leq (T_0, T_1)$  if and only if  $T'_0 \subseteq T_0$  and  $T'_1 \subseteq T_1$ . A pair  $(P_0, P_1)$  of subsets of  $M$  is a *pair of paths through*  $(T_0, T_1)$  if  $P_i$  is a path through  $T_i$  ( $i = 0, 1$ ). Then the notions of a *dense* set and a *generic* path can be defined for  $(\mathcal{T}_S^2, \leq)$  in the same way as for  $(\mathcal{T}_S, \subseteq)$ .

**Lemma 16.** *Let  $(M, S)$  be a countable model of  $\mathbf{RCA}_0$ . For each tree  $T \in \mathcal{T}_S$ , there exist two paths  $P_0$  and  $P_1$  through  $T$  such that  $P_i$  is not in  $\Delta_1^0(S \cup \{P_{1-i}\})$  ( $i = 0, 1$ ), and  $(M, S \cup \{P_0, P_1\}) \models \Sigma_1^0$  induction.*

*Proof.* Let  $(M, S)$  be a countable model of  $\mathbf{RCA}_0$ . Let  $\{\exists y \theta_e(x, X[y]) : \theta_e(x, s) \in \Sigma_0^0, e \in \omega\}$  enumerate all the  $\Sigma_1^0$ -formulas with parameters from  $M \cup S$  and no free variables other than  $x$  and  $X$ . For each  $(e, d) \in \omega^2$ , we define  $D_{e,d}^1$  to be the set of all  $(T_0, T_1) \in \mathcal{T}_S^2$  such that  $(M, S)$  satisfies, for some  $m \in M$ , either

- A1.  $\forall s \in T_1[\text{lh}(s) = m + 1 \rightarrow s(m) = 1] \wedge$ 
  - (i)  $\forall t \in T_0 \neg \theta_e(m, t) \vee$
  - (ii)  $\exists w \forall t \in T_0[\text{lh}(t) = w \rightarrow \exists y < w \theta_d(m, t[y])]$  or
- A2.  $\forall s \in T_1[\text{lh}(s) = m + 1 \rightarrow s(m) = 0] \wedge$ 
  - (i)  $\forall t \in T_0 \neg \theta_d(m, t) \vee$

(ii)  $\exists w \forall t \in T_0 [\text{lh}(t) = w \rightarrow \exists y < w \theta_e(m, t[y])]$ .

To show that  $D_{e,d}^1$ 's are dense, fix any  $(T_0, T_1) \in \mathcal{T}_S^2$ . Let  $P_1 \subseteq M$  be a path through  $T_1$ . Since  $P_1 \notin S$ , by Lemma 8  $T_0$  has a path  $P_0$  such that  $P_1$  is not in  $\Delta_1^0(S \cup \{P_0\})$ . Then there exists  $m \in M$  such that one of the following conditions holds:

- B1.  $m \in P_1 \wedge (M, S \cup \{P_0\}) \models \forall y \neg \theta_e(m, P_0[y])$ ,
- B2.  $m \in P_1 \wedge (M, S \cup \{P_0\}) \models \exists y \theta_d(m, P_0[y])$ ,
- B3.  $m \notin P_1 \wedge (M, S \cup \{P_0\}) \models \exists y \theta_e(m, P_0[y])$ ,
- B4.  $m \notin P_1 \wedge (M, S \cup \{P_0\}) \models \forall y \neg \theta_d(m, P_0[y])$ .

Suppose that condition B1 holds. Let  $t_1 = P_1[m+1]$  and  $T'_1 = \{s \in T_1 : s \subseteq t_1 \text{ or } t_1 \subseteq s\}$ . Let  $T'_0 = \{s \in T_0 : \forall t \subseteq s \neg \theta_e(m, t)\}$ . Then it is clear that  $(T'_0, T'_1) \in \mathcal{T}_S^2$  and A1.(i) holds with  $(T'_0, T'_1)$ . Thus  $(T'_0, T'_1) \in D_{e,d}^1$ . Suppose that condition B2 holds. Take  $T'_1$  as above. Let  $t_0$  be an initial segment of  $P_0$  with  $\theta_d(m, t_0)$ . Let  $T'_0 = \{s \in T_0 : s \subseteq t_0 \text{ or } t_0 \subseteq s\}$ . Then  $(T'_0, T'_1) \in \mathcal{T}_S^2$  satisfies A1.(ii). Thus  $(T'_0, T'_1) \in D_{e,d}^1$ . The other two cases can be treated similarly. Hence  $D_{e,d}^1$ 's are dense.

We also define  $D_{e,d}^2$  in the same way as  $D_{e,d}^1$  but exchanging the roles of  $T_0$  and  $T_1$ . Then  $D_{e,d}^2$ 's are also dense.

For  $e \in \omega$ ,  $m \in M$ , we define  $D_{m,e}^3$  to be the set of  $(T_0, T_1) \in \mathcal{T}_S^2$  such that  $(M, S)$  satisfies  $\forall n \leq m$  [

- 1.  $\exists w \forall s_0 \in T_0 \forall s_1 \in T_1 [\text{lh}(s_0) = \text{lh}(s_1) = w \rightarrow (\exists y < w \theta_e(n, s_0 \oplus s_1[y]))]$  or
- 2.  $\forall s_0 \in T_0 \forall s_1 \in T_1 \neg \theta_e(n, s_0 \oplus s_1)$ ,

where  $s_0 \oplus s_1$  is the 0-1 sequence  $t$  with length  $\text{lh}(s_0) + \text{lh}(s_1)$  defined by  $t(2n) = s_0(n)$  and  $t(2n+1) = s_1(n)$ . By the proof of Lemma 3, we can easily see that  $D_{m,e}^3$ 's are dense.

We now have a countable set  $\mathcal{D} = \{D_{e,d}^1, D_{e,d}^2, D_{m,e}^3 : m \in M, e, d \in \omega\}$  of dense sets. Let  $(P_0, P_1)$  be a pair of paths through  $(T, T)$  which is generic for  $\mathcal{D}$ . By the proof of Lemma 3, it is easy to see  $(M, S \cup \{P_0, P_1\}) \models \Sigma_1^0$  induction since  $(P_0, P_1)$  is generic for  $\{D_{m,e}^3\}$ . Since  $(P_0, P_1)$  is generic for  $\{D_{e,d}^1\}$ , there exists  $m \in P_1$  such that  $\forall y \neg \theta_e(m, P_0[y]) \vee \exists y \theta_d(m, P_0[y])$ , or  $m \notin P_1$  such that  $\exists y \theta_e(m, P_0[y]) \vee \forall y \neg \theta_d(m, P_0[y])$ . Thus it is not the case that

$$P_1 = \{m : \exists y \theta_e(x, P_0[y])\} = \{m : \forall y \neg \theta_d(x, P_0[y])\}.$$

So  $P_1$  is not in  $\Delta_1^0(S \cup \{P_0\})$ . In the similarly way, we can prove that  $P_0 \notin \Delta_1^0(S \cup \{P_1\})$ . This completes the proof.  $\square$

To prove our main theorem, we generalize Lemma 16 as follows.

**Lemma 17.** *Let  $(M, S)$  be a countable model of  $\text{RCA}_0$ . For any tree  $T \in \mathcal{T}_S$ , there exist two paths  $P_0$  and  $P_1$  through  $T$  such that  $\{(P_0)_i : i \in M\} \cap \{(P_1)_j : j \in M\} \subseteq S$  and  $(M, S \cup \{P_0, P_1\}) \models \Sigma_1^0$  induction.*

*Proof.* Let  $(M, S)$  be a countable model of  $\text{RCA}_0$ . Let  $\{\exists y \theta_e(x, X[y]) : \theta_e(x, s) \in \Sigma_0^0, e \in \omega\}$  enumerate all the  $\Sigma_1^0$ -formulas with parameters from  $M \cup S$  and no free variables other than  $x$  and  $X$ . Let  $D_{j,e,d}^4$  be the set of  $(T_0, T_1) \in \mathcal{T}_S^2$  such that  $(Z)_j \in S$  for any path  $Z$  through  $T_1$ , or that there exists  $m \in M$  such that  $(M, S)$  satisfies that

- A1.  $\forall s \in T_1 [\text{lh}(s) = \langle m, j \rangle + 1 \rightarrow s(\langle m, j \rangle) = 1] \wedge$ 
  - (i)  $\forall t \in T_0 \neg \theta_e(m, t) \vee$
  - (ii)  $\exists w \forall t \in T_0 [\text{lh}(t) = w \rightarrow \exists y < w \theta_d(m, t[y])]$  or

- A2.  $\forall s \in T_1[\text{lh}(s) = \langle m, j \rangle + 1 \rightarrow s(\langle m, j \rangle) = 0] \wedge$   
 (i)  $\forall t \in T_0 \neg \theta_d(m, t) \vee$   
 (ii)  $\exists w \forall t \in T_0[\text{lh}(t) = w \rightarrow \exists y < w \theta_e(m, t[y])].$

By straightforward adaptation of the proof of Lemma 16,  $D_{j,e,d}^4$ 's are dense.

Let  $D_{m,e}^3$ 's be the dense sets defined in the proof of Lemma 16 and let  $\mathcal{D} = \{D_{m,e}^3, D_{j,e,d}^4 : e, d \in \omega, j, m \in M\}$ . Then  $\mathcal{D}$  is a countable set of dense sets. So we can take a pair  $(P_0, P_1)$  of paths through  $(T, T)$  which is a generic for  $\mathcal{D}$ . It is easy to see that  $(M, S \cup \{P_0, P_1\}) \models \Sigma_1^0$  induction. Thus we only need to show that  $\{(P_0)_i : i \in M\} \cap \{(P_1)_j : j \in M\} \subseteq S$ .

To see this, we first claim that for each  $j \in M$ , either  $(P_1)_j \in S$  or  $(P_1)_j \notin \Delta_1^0(S \cup \{P_0\})$ . Fix any  $j \in M$  such that  $(P_1)_j \notin S$ . Since  $(P_0, P_1)$  is generic for  $\{D_{j,e,d}^4\}$  and  $(P_1)_j \notin S$ , there exists  $m \in M$  such that  $(M, S)$  either  $(P_1)_j \neq \{m : (M, S \cup \{P_0\}) \models \exists y \theta_e(m, P_0[y])\}$  or  $(P_1)_j \neq \{m : (M, S \cup \{P_0\}) \models \forall y \neg \theta_d(m, P_0[y])\}$ . Since  $(e, d)$  is arbitrary,  $(P_1)_j$  is not in  $\Delta_1^0(S \cup \{P_0\})$ .

Choose any  $X \in \{(P_0)_i : i \in M\} \cap \{(P_1)_j : j \in M\}$ . Then there exist  $i, j$  such that  $X = (P_0)_i = (P_1)_j$ . The above claim implies  $X \in S$  since  $(P_1)_j = (P_0)_i \in \Delta_1^0(S \cup \{P_0\})$ . Hence  $\{(P_0)_i : i \in M\} \cap \{(P_1)_j : j \in M\} \subseteq S$ . This completes the proof.  $\square$

**Lemma 18.** *Let  $M$  be a countable model of  $\mathbf{IS}_1$ . Let  $A$  be a subset of  $M$  such that  $(M, \{A\})$  satisfies  $\Sigma_1^0$  induction. Then there exist three sets  $S_1, S_2, S_3$  of subsets of  $M$  such that*

- (1)  $(M, S_i) \models \text{WKL}_0$  ( $i = 1, 2, 3$ ), and
- (2)  $S_1 \cap S_2 = \Delta_1^0(\{A\})$ , and
- (3)  $S_1 \cup S_2 \subseteq S_3$ .

*Proof.* By Theorem 7, there is a  $\Pi_1^0$  formula  $\psi(X, Y)$  such that for any model  $(M, S)$  of  $\text{WKL}_0$  and for any  $B \in S$ , (i) there exists  $W \in S$  such that  $(M, S) \models \psi(B, W)$ , and (ii) if  $(M, S) \models \psi(B, W)$ , then  $(M, \{(W)_n : n \in M\}) \models \text{WKL}_0$  and  $B \in \{(W)_n : n \in M\}$ . Let  $(M, \{A\})$  be a countable model of the axioms of ordered semirings and  $\Sigma_1^0$  induction, and let  $S_0 = \Delta_1^0(\{A\})$ . By the normal form theorem, write  $\psi(X, Y)$  as  $\forall n \theta(X[n], Y[n])$  where  $\theta$  is  $\Sigma_0^0$ . Let  $T$  be the set of  $\tau \in 2^{<M}$  such that  $\forall \sigma \subseteq \tau \theta(A[\text{lh}(\sigma)], \sigma)$ . Then  $T \in S_0$  and by (i),  $T$  is infinite. We claim that  $T$  has no path in  $S_0$ . To see this, suppose that  $T$  has a path  $P$  in  $S_0$ . By Lemma 5, let  $(M, S')$  be a model of  $\text{WKL}_0$  with  $S_0 \subseteq S'$ . By (ii),  $(M, \{(P)_n : n \in M\}) \models \text{WKL}_0$ . But  $(M, S_0)$  is the least  $M$ -model of  $\text{RCA}_0$  containing  $A$ . This is a contradiction. Thus  $T$  has no path in  $S_0$ .

By Lemma 17, there exist two paths  $P_1, P_2$  through  $T$  such that  $\{(P_1)_i : i \in M\} \cap \{(P_2)_j : j \in M\} \subseteq S_0$  and  $(M, \{A, P_1, P_2\}) \models \Sigma_1^0$  induction. By Lemma 5, we have a model  $(M, S_3)$  of  $\text{WKL}_0$  containing  $A, P_1$  and  $P_2$ . Let  $S_1 = \{(P_1)_i : i \in M\}$  and  $S_2 = \{(P_2)_j : j \in M\}$ . Then  $(M, S_1)$  and  $(M, S_2)$  are models of  $\text{WKL}_0$  such that  $S_1 \cap S_2 = S_0$  and  $S_1 \cup S_2 \subseteq S_3$ .  $\square$

**Theorem 19.** *Let  $\varphi(x, X, Y)$  be an arithmetical formula with exactly the free variables shown. If  $\text{WKL}_0$  proves  $\forall x \forall X \exists! Y \varphi(x, X, Y)$ , then so does  $\text{RCA}_0$ .*

*Proof.* Let  $\varphi(x, X, Y)$  be an arithmetical formula with exactly the free variables shown. Suppose that  $\text{WKL}_0$  proves  $\forall x \forall X \exists! Y \varphi(x, X, Y)$  and  $\text{RCA}_0$  can not prove it. Then by Gödel's completeness theorem, there exists a countable model  $(M, S)$  of  $\text{RCA}_0$  such that  $\neg \exists! Y \varphi(n, A, Y)$  holds in  $(M, S)$  for some  $n \in M$  and  $A \in S$ .

First suppose that  $\exists Y\varphi(n, A, Y)$  holds in  $(M, S)$ . Then there exists more than one set in  $S$  which satisfies  $\varphi$ . By Lemma 5, there exists a model  $(M, S')$  of  $\text{WKL}_0$  such that  $S \subseteq S'$ . Obviously,  $S'$  has at least two distinct sets which satisfy  $\varphi$ . Hence  $\text{WKL}_0$  does not prove  $\forall x\forall X\exists!Y\varphi(x, X, Y)$ , which is a contradiction.

Next assume that  $\forall Y\neg\varphi(n, A, Y)$  holds within  $(M, S)$ . Let  $S_0 = \Delta_1^0(\{A\})$ . Then  $\forall Y\neg\varphi(n, A, Y)$  holds within  $(M, S_0)$ . By Lemma 18, there exist three sets  $S_1, S_2, S_3$  of subsets of  $M$  such that (1)  $(M, S_i) \models \text{WKL}_0$  ( $i=1, 2, 3$ ); (2)  $S_1 \cap S_2 = S_0$ ; (3)  $S_1 \cup S_2 \subseteq S_3$ . By (1), there exists  $B_i \in S_i$  such that  $(M, S_i) \models \varphi(n, A, B_i)$  for  $i = 1, 2$ . Since there exists no set  $C$  in  $S_0$  which satisfies  $\varphi(n, A, C)$ , we must have  $B_1 \neq B_2$  by (2). So,  $(M, S_3) \models \neg\exists!Y\varphi(n, A, Y)$  by (3). Since  $(M, S_3) \models \text{WKL}_0$ , this contradicts with the assumption that  $\text{WKL}_0$  proves  $\forall x\forall X\exists!Y\varphi(x, X, Y)$ . Then the proof is completed.  $\square$

**Open problems.** (1) If  $\text{WKL}_0$  proves  $\forall x\forall X\exists!Y\varphi(x, X, Y)$  and  $\varphi$  is  $\Pi_1^1$ , then does  $\text{RCA}_0$  already prove it?

(2) If  $\text{WKL}_0 + \Sigma_n^0$  induction proves  $\forall x\forall X\exists!Y\varphi(x, X, Y)$  and  $\varphi$  is arithmetical, then does  $\text{RCA}_0 + \Sigma_n^0$  induction already prove it ( $n = 2, 3, \dots$ )? (cf. [1].)

We remark that (2)'s answer is known to be affirmative when  $n = \infty$  (i.e., full arithmetical induction). To see this, let  $(M, S)$  be a model of  $\text{ACA}_0$ . By a formalized version of Kreisel's hard core theorem which is provable in  $\text{ACA}_0$  (Corollary VIII.2.26 [8]), for each  $A \in S$ , there exist  $W, W' \in S$  such that  $(M, \{(W)_n\})$  and  $(M, \{(W')_n\})$  are models of  $\text{WKL}_0$  and  $\{(W)_n\} \cap \{(W')_n\} = \Delta_1^0(\{A\})$ . The rest of the argument is the same as the proof of Theorem 19.

We also remark that Fernandes [2] solves the problem (2) with restriction to the case that  $n = 2$  and  $\varphi$  is  $\Sigma_3^0$ . For more related problems, see [10].

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MATHEMATICAL INSTITUTE, TOHOKU UNIVERSITY SENDAI 980-8578, JAPAN  
*E-mail address:* tanaka@math.tohoku.ac.jp, 95m41@math.tohoku.ac.jp